

CSCD 433

Network Programming

Winter 2017



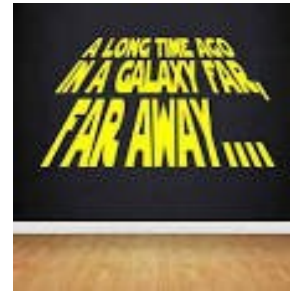
Lecture 7

Ethernet and Wireless 802.11

Topics

- 802 Standard
 - MAC and LLC Sublayers
 - Review of MAC in Ethernet
 - Ethernet vs. Wireless - CSMA
 - MAC in 802.11 Wireless

IEEE Standards



- In 1985, Computer Society of IEEE started a project, called Project 802,
- Set standards to enable intercommunication among equipment from a variety of manufacturers ... wonderful!!!
- Project 802
 - Specifies functions of physical layer and the data link layer of major LAN protocols

IEEE 802 Series of LAN Standards

802 standards free to
download from

[http://standards.ieee.org
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IEEE 802®: Overview & Architecture

IEEE 802.1™ Bridging & Management

IEEE 802.2™: Logical Link Control

IEEE 802.3™: CSMA/CD Access Method

IEEE 802.4™: Token-Passing Bus Access Method

IEEE 802.5™: Token Ring Access Method

IEEE 802.6™: DQDB Access Method

IEEE 802.7™: Broadband LAN

IEEE 802.10™: Security

IEEE 802.11™: Wireless

IEEE 802.12™: Demand Priority Access

IEEE 802.16™: Broadband Wireless Metropolitan Area Networks

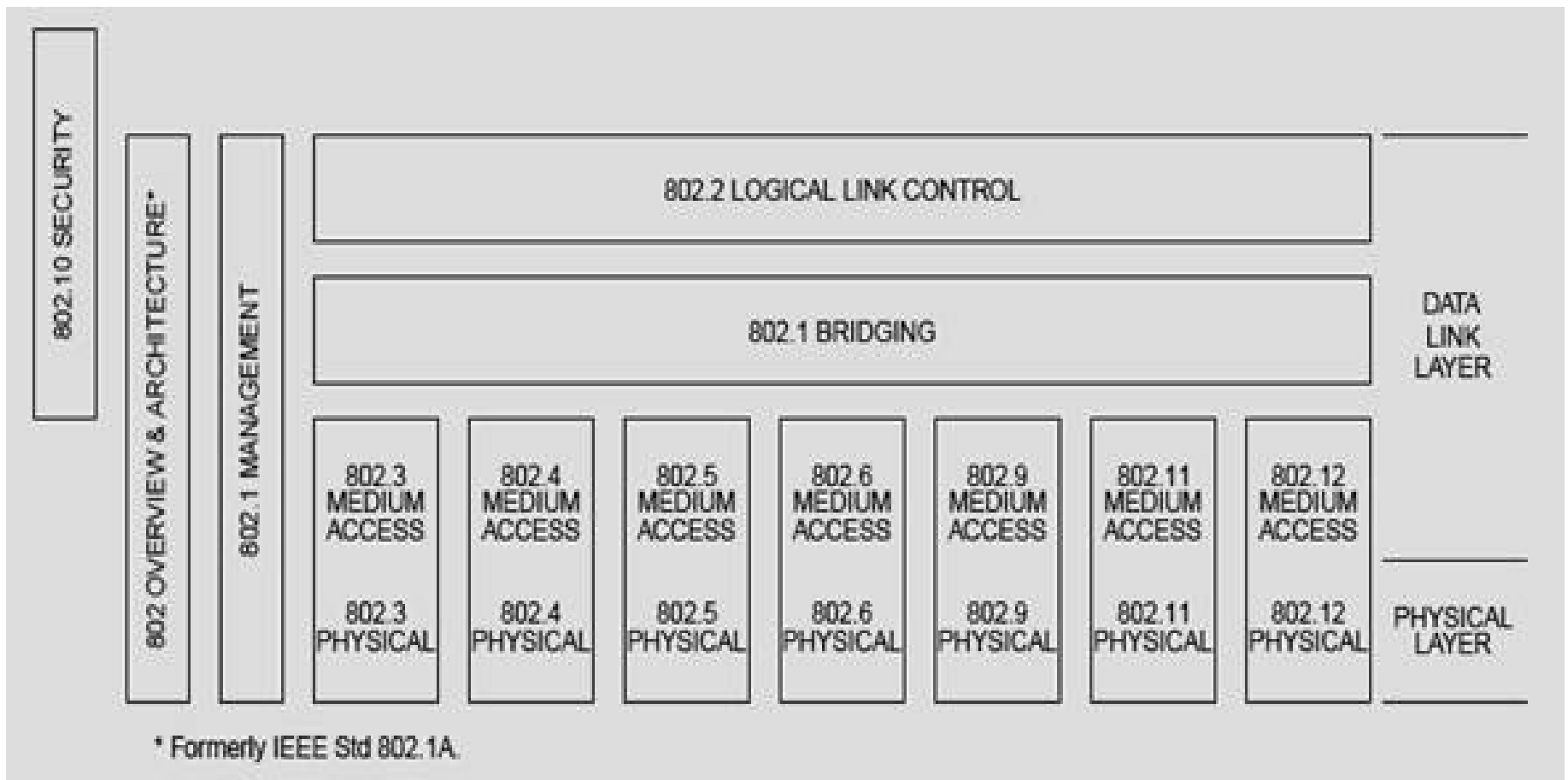
WiMAX

802.11 Physical Layer

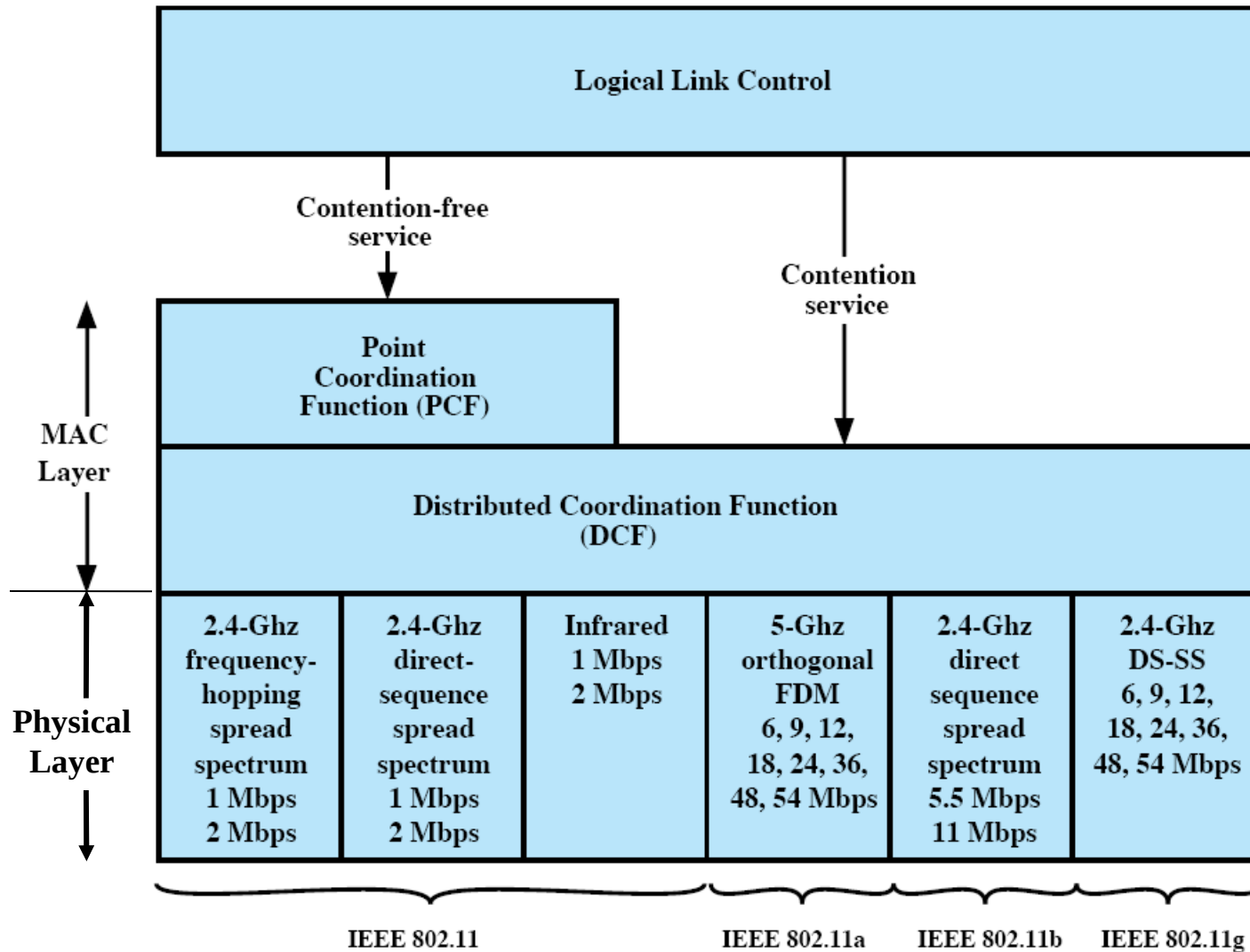
- Issued in four stages
- **1997, First part**
 - IEEE 802.11
 - Includes MAC layer and three physical layer specifications
 - Two in 2.4-GHz band and one infrared, Operate 1 and 2 Mbps
- **1999, Two more parts**
 - IEEE 802.11a
 - 5-GHz band, data rate up to 54 Mbps
 - IEEE 802.11b
 - 2.4-GHz band, data rate at 5.5 and 11 Mbps
- **2002, Most recent**
 - IEEE 802.11g extended IEEE 802.11b to higher data rates, up to 54 Mbps
- **At present**
 - IEEE 802.11n and 802.11ac data rate up to hundreds of Mbps

IEEE 802 Standard

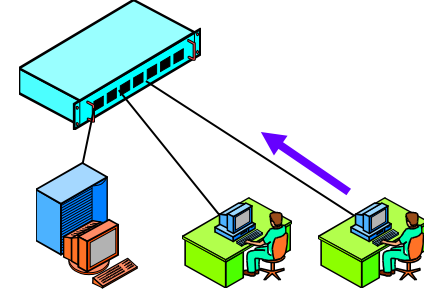
Complete specification of 802 standard



IEEE 802.11 Protocol Architecture



802 Layering



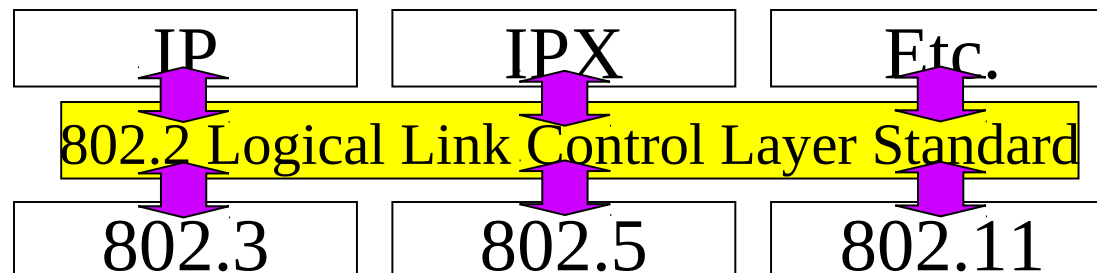
Media Access Control (MAC) Layer

- 802.11 standard specifies common medium access control (MAC) Layer
- In general, MAC Layer manages and maintains communications between 802.11 stations by coordinating access to shared radio channel, has most of functionality

Logical Link Control (LLC) Layer

Adds optional error correction (rarely used)

Connects to next-higher-layer (internet), multiplexes higher level protocols



Medium Access Control

- Two sublayers

Lower sublayer

Distributed Coordination Function (DCF)

- Uses a contention algorithm to provide access to all traffic
- Uses CSMA/CA – Carrier Sense Multiple Access/Collision Avoidance

Higher sublayer

Point Coordination Function (PCF)

- Uses a centralized algorithm
 - “Contention free”
 - Implemented on top of DCF
- **Note:** PCF has not been popularly implemented in today’s 802.11 products, DCF is widely used

MAC Sublayer Functions

- **MAC** is set of rules to determine how to access medium and data link components.

It provides core framing operations, sets the frame header fields

- **MAC purpose in General**

- Coordinates and shares use of radio bandwidth
- Synchronization between stations
- Datagram transfer function

- **MAC layer management functions**

Authentication and De-Authentication

Association, Re-Association, and Disassociation

Beacon and Probe frames

MAC Sublayer Functions

- 802.11 uses CSMA/CA mechanism

Carrier Sense Multiple Access with Collision Avoidance

- It is considered to be 'fair' for all users because treats them equally

Recall that Ethernet uses CSMA/CD



Brief Review of Classical or Standard Ethernet

Review of Classical Ethernet

- Recall that classical Ethernet is shared technology
- Everyone has access to wires
- Users contend with collisions and MAC layer protocol dealt with these collisions
 - Note – This is with traditional cable and Hubs
- Review characteristics of Ethernet

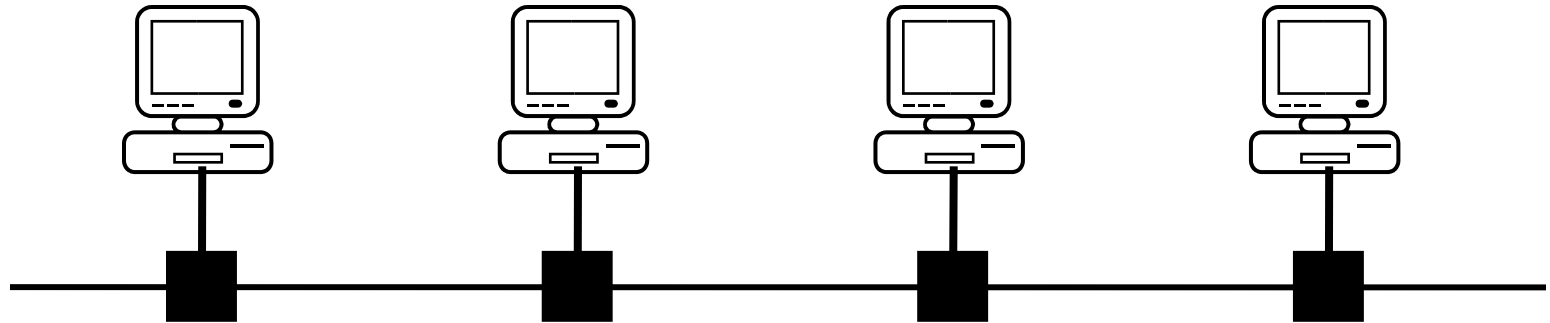
Ethernet Recap



- **Classic Ethernet**
 - One long cable, 500 meter max segment
 - Snaked around building as single, long cable
 - All computers attached
- **Thick Ethernet and Thin**
 - Began as thick yellow cable, marked every 2.5 meters to show computer attachments
 - Thinner, bent more easily connections with BNC connectors
 - Cheaper to install, 185 meter max segment



CSMA/CD Protocol



All hosts transmit & receive on one channel

When a host has a packet to transmit

1. **Carrier Sense:** Check that the line is quiet before transmitting
2. **Collision Detection:** Detect collision as soon as Possible.
Collision is detected, stop transmitting; wait a **random time**, then return to step 1.

Ethernet CSMA/CD algorithm

Algorithm

1. NIC receives datagram from network layer, creates frame
2. If NIC senses channel idle, starts frame transmission
If NIC senses channel busy, waits until channel idle, then transmits
3. If NIC transmits entire frame without detecting another transmission, NIC is done with frame !

Ethernet CSMA/CD Algorithm

4. If NIC detects another transmission while transmitting, aborts and sends jam signal

5. **After aborting**

NIC enters **exponential backoff**

after **mth** collision, NIC chooses a K , small integer, at random from $\{0, 1, 2, \dots, 2^m - 1\}$

NIC then waits $K \cdot 512$ bit time,

- Returns to Step 2

Ethernet's CSMA/CD (more)

Exponential Backoff

- **Goal** Adapt retransmission attempts to estimated current load
 - **Heavy load -> random wait will be longer and more varied**
- **First collision**: Choose K from $\{0,1\}$;
Delay is $K \cdot 512$ bit transmission times
- **After second collision**: Choose K from $\{0,1,2,3\} \dots$
- **After ten collisions**, Choose K from $\{0,1,2,3,4,\dots,1023\}$
- Set size grows **Exponentially**

Wireless Communication Systems

- In terms of packet or frame delivery
- What complicates wireless networking vs. wired networking?

Wireless Link Characteristics



Differences from wired link

- **Decreased signal strength:** Radio signal attenuates as it propagates through matter (path loss)
- **Interference from other sources:** Standardized wireless network frequencies shared by other devices that can interfere
- **Multipath propagation:** Radio signal reflects off objects and ground, arriving at destination at slightly different times

.... make communication across link much more “difficult”

CSMA/CD vs. CSMA/CA

- **CSMA/CD – CSMA/Collision detection**
 - For wired communication
 - No control BEFORE transmission
 - Generates collisions
 - Collision Detection- Monitors signal strength, can detect
- **CSMA/CA – CSMA/Collision Avoidance**
 - For wireless communication
 - Collision avoidance BEFORE transmission
 - Difference in energy/power for transmit & receive
 - Difficult to distinguish between incoming weak signals, noise, and effects of own transmission

802.11 Medium Access Control

- Shares the medium through coordination with other stations
 1. Sends control packets for coordination
Including Acks
 2. Sets and sends individual frame timers for all to see, each frame has its own timer
- Note: There is the 802.11 standard and then there is reality of did it actually get implemented!
 - Create a lot of details for a standard but not all of it is implemented

Wireless Collision Avoidance

STEPS

1. Have a frame to send
2. Wait a random time, until channel is idle
3. Sense it is idle for short time, called **DIFS** period
4. Sends frame, if gets through, destination waits a **SIFS** time and sends an **ACK**
5. Lack of an ACK back means frame failed
6. Sender then doubles backoff time, tries again
7. Continues until frame succeeds

Distributed Inter-frame Spacing (DIFS)
Short Inter-frame Spacing (SIFS)

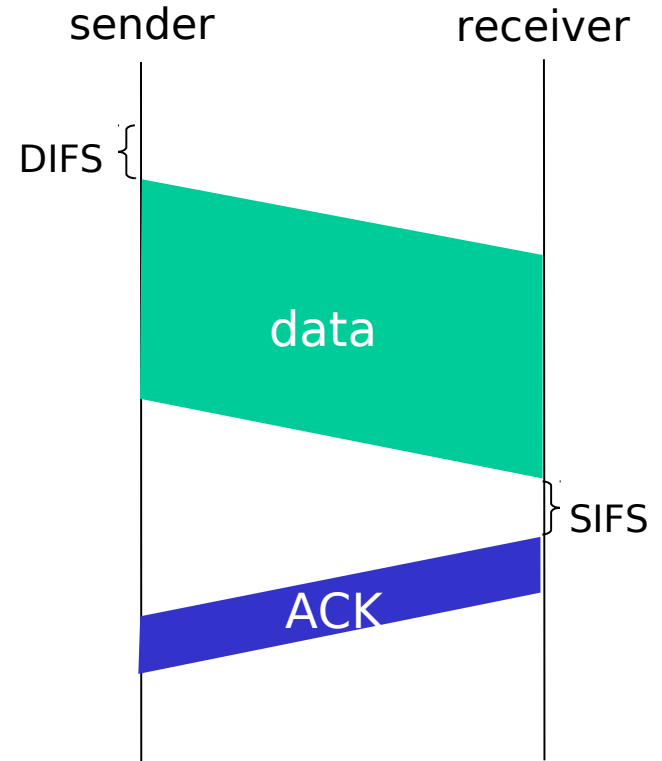
IEEE 802.11 MAC Protocol: CSMA/CA

802.11 sender

1. If sense channel idle for **DIFS** then transmit entire frame
2. If sense channel busy then
 - a) start random backoff time
 - b) timer counts down
 - c) transmit when timer expires
 - d) if no ACK, increase random backoff interval, repeat 2 (frame failed)

802.11 receiver

- if frame received OK
return ACK after **SIFS** (ACK needed due to hidden terminal problem)



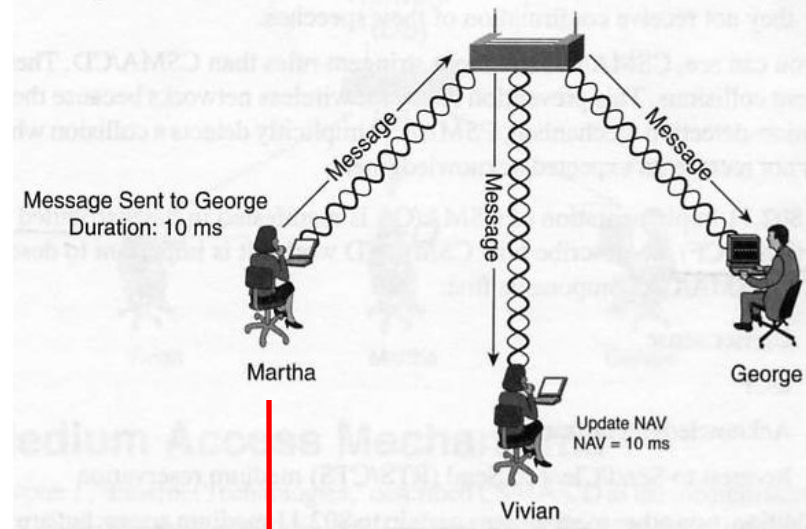
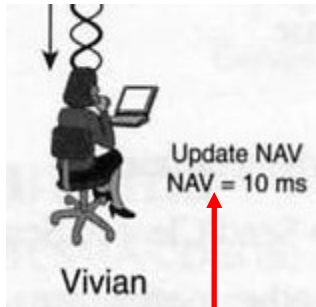
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Short Inter-frame Spacing (SIFS)

802.11 Frames have Timers

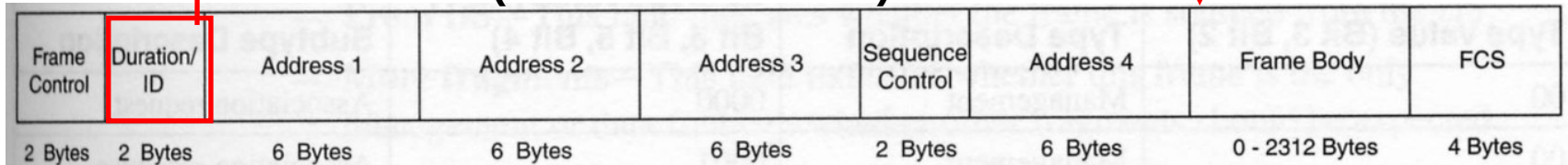
- **Network Allocation Vector (NAV) timer**
 - NAV is set when a frame sequence is sent
 - Says how long a sequence will take so other stations have an idea when the medium will be available
 - For example, a NAV for a data frame will also include the ACK back

Next Slides Demo this with Example ...

NAV Timer

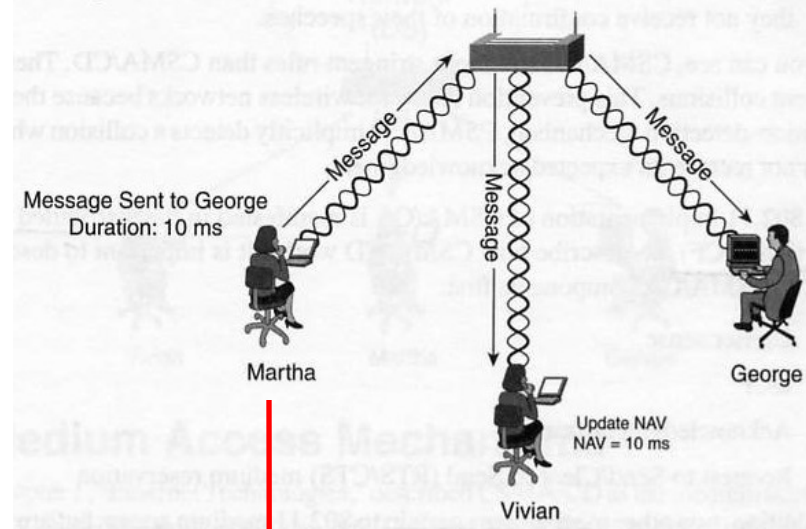
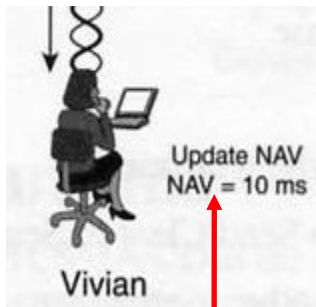


General 802.11 Frame (more on this later)

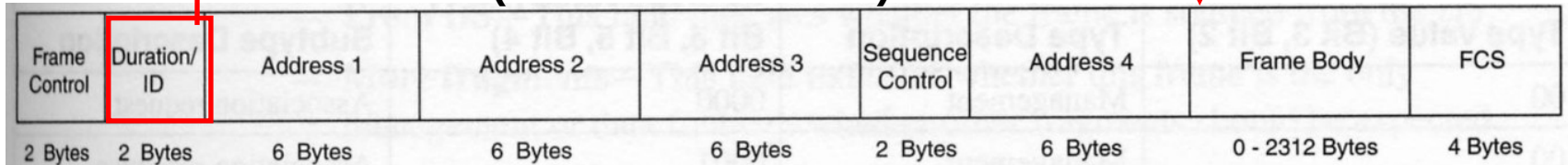


- All stations have a **NAV (Network Allocation Vector) timer**
Protects frames from interruption
- **Example: Martha sends a frame to George**
- Since wireless medium is “broadcast-based” shared medium, all stations including Vivian receive frame
- Vivian updates her NAV timer with duration value
- Vivian will not attempt to transmit until her NAV is decremented to 0.
- Stations will only update their NAV when duration field value received is greater than their current NAV

Duration Field



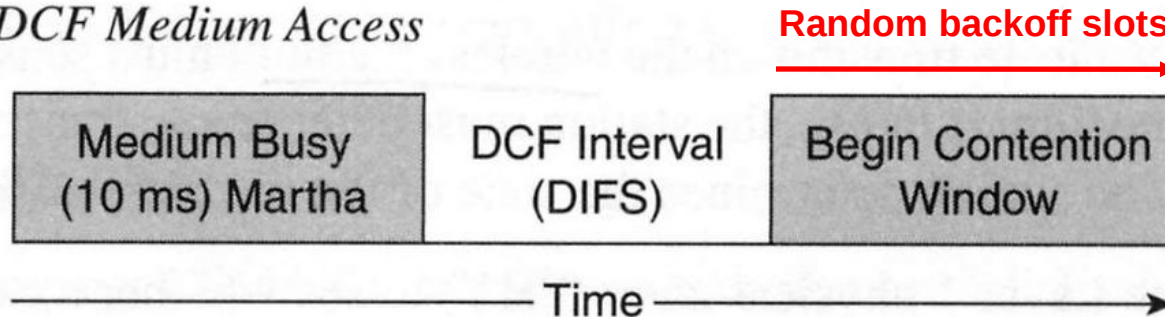
General 802.11 Frame (more on this later)



- **Duration/ID field** – Number of microseconds medium is expected to remain busy for transmission currently in progress
 - **Transmitting device sets Duration time in microseconds**
 - **Includes time to:**
 - **Transmit this frame to AP (or to the client if an AP)**
 - **Includes returning ACK**
- All stations monitor this field!
- All stations update their **NAV** (Network Allocation Vector) timer

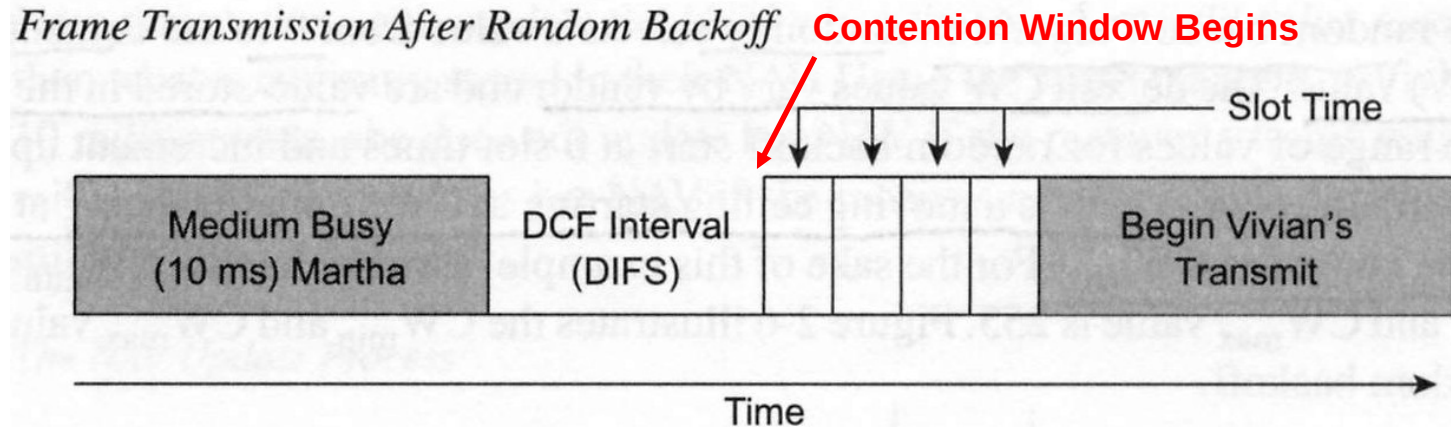
Wanting to transmit (1/3)

Timeline for DCF Medium Access



- Station wanting to transmit.
- **Carrier Sensing**
 - **Physical: Physically senses medium is idle**
 - **Virtual: NAV timer is 0**
- Waits **DIFS** (Distributed or DCF Interframe Space)
 - **Minimum amount of medium idle time until contention-based services begin.**
 - **Once DIFS is over, stations can contend for access.**
- **Contention window begins.**
 - **Uses random backoff algorithm to determine when it can attempt to access the medium. (next)**

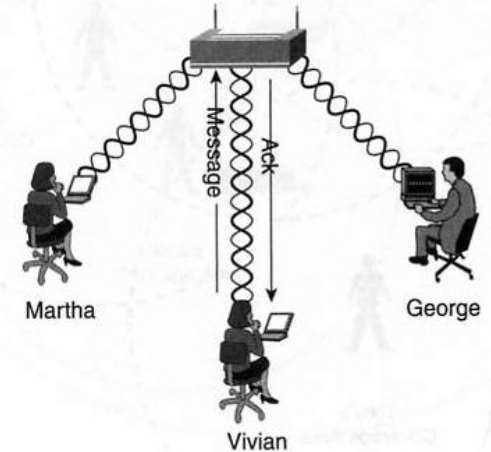
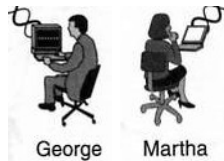
Wanting to transmit (2/3)



- The **random backoff algorithm** randomly selects a value from **0 to 255**. Maximum value varies by vendor.
- The random value is number of **802.11 slot times** the station must wait after the DIFS, during the contention window before it may transmit.
- **Stations pick a random slot** and wait for that slot before attempting to access the medium.
- With several stations attempting to transmit, the station that picks the **lowest slot, lowest random number, wins**.

Wanting to transmit (3/3)

Others
update NAV



General 802.11 Frame (more on this later)

Frame Control	Duration/ID	Address 1	Address 2	Address 3	Sequence Control	Address 4	Frame Body	FCS
2 Bytes	2 Bytes	6 Bytes	6 Bytes	6 Bytes	2 Bytes	6 Bytes	0 - 2312 Bytes	4 Bytes

- **Station transmits**, setting the **Duration ID** to the time needed to transmit data, ACK and IFSSs.
- **Other stations** with higher slots will see the new transmission and wait to transmit.
- If **frame arrives at AP** (assuming the transmitter is a station), then an **ACK will be returned**
- If there is **not an ACK received**, the sending station assumes there has been a collision
 - **If two stations have the same lowest slot time and both transmit, then a collision occurs**
- Stations will **update its retry counter** (double) to determine a **new randomly selected slot time** and **process starts all over again**

Hidden Terminal Problem in WLANs

- Both H1 and H2 transmit at same time
- Signals collide at AP, H1 can't detect H2

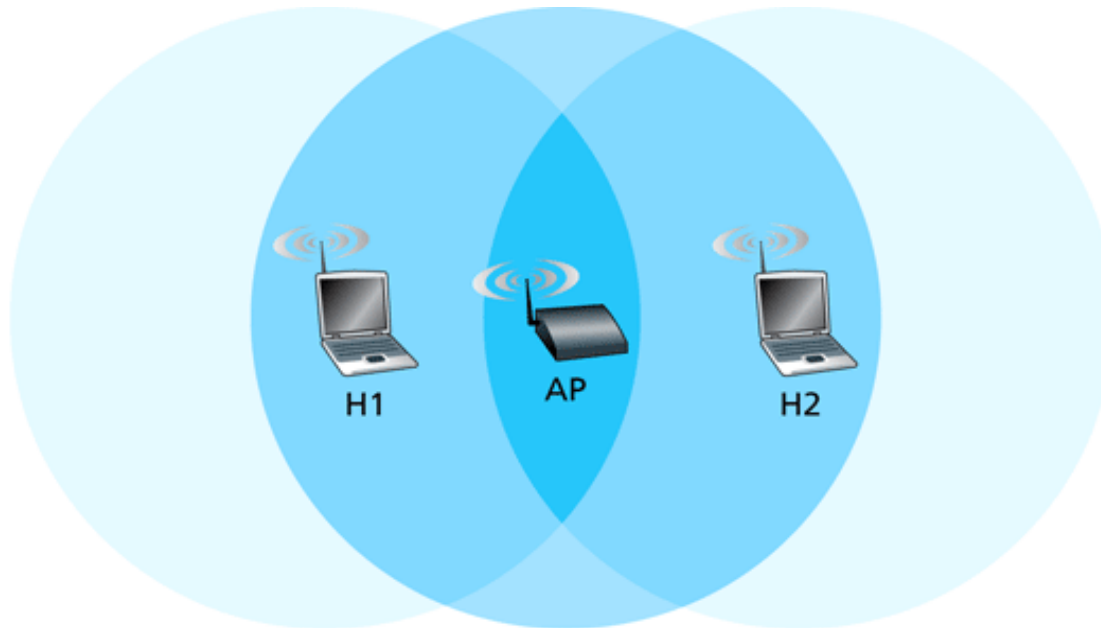


Figure 6.11 ♦ Hidden terminal example: H1 is hidden from H2, and vice versa.

Avoiding collisions: RTS/CTS

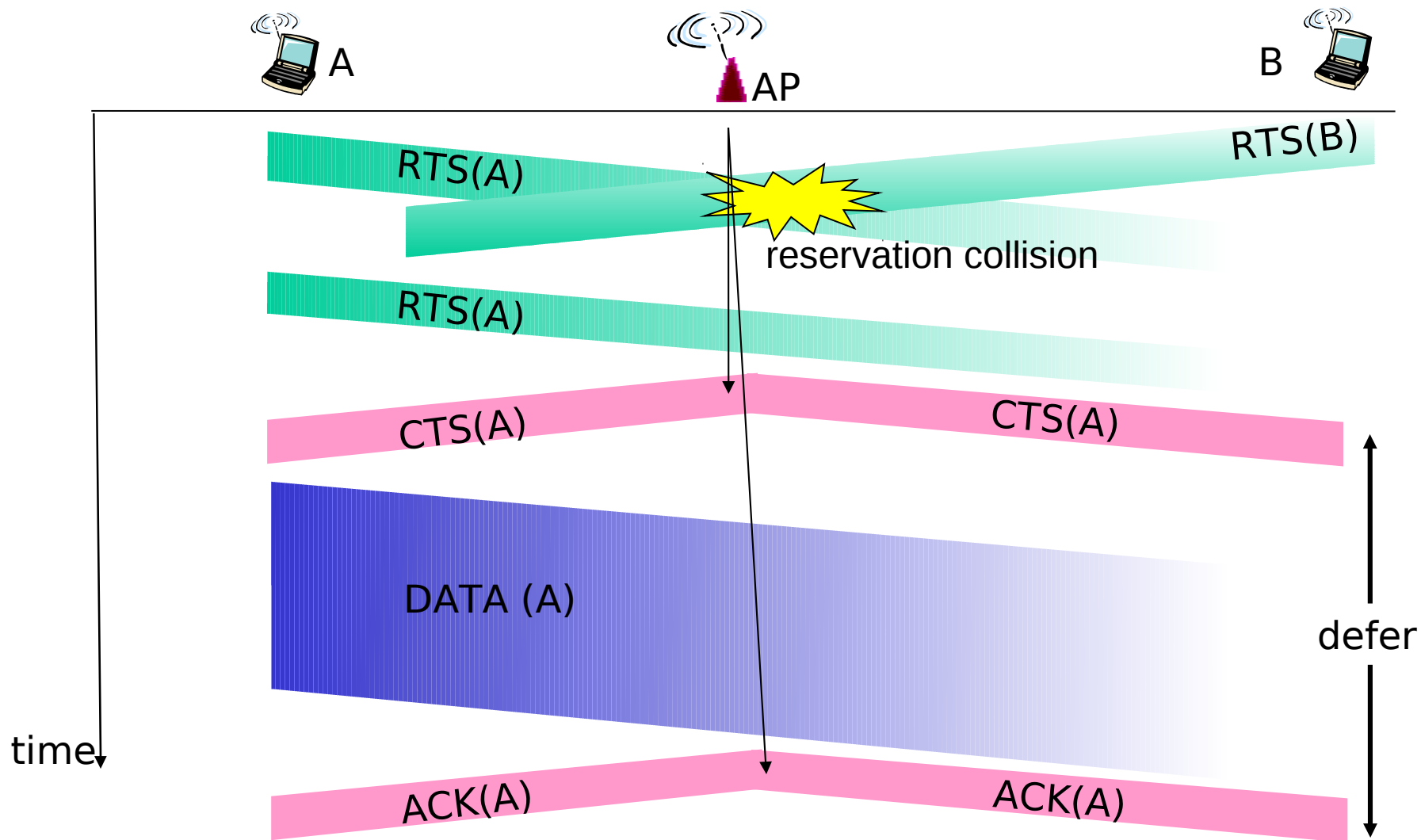
Another Idea: Allow sender to “reserve” channel rather than random access of data frames: avoid collisions of long data frames

Sender first transmits *small* request-to-send (RTS) packets using CSMA, say want to transmit to AP

- RTSs may still collide with each other (but they're short)
- AP broadcasts clear-to-send (CTS) in response to RTS
- RTS heard by all nodes
 - Sender transmits data frame
 - Other stations defer transmissions

Avoids data frame collisions completely
using small reservation packets!

Collision Avoidance: RTS-CTS exchange



RTS/CTS in practice

- 802.11 standardized both CSMA/CA and RTS/CTS
- In practice, most operators **disable** RTS/CTS
 - **Very high overhead!**
 - RTS/CTS packets sent at “base rate” (often 1Mbit)
 - Neighboring AP's are often configured to use non-overlapping channels, so hidden terminals on downlink are rare



MAC Addresses

Review

MAC Addresses

Network Layer

32-bit IP address

Network-layer address, dotted decimal

Ex.: 146.187.130.76

To route datagram to destination machine, router uses

MAC (or LAN or physical or Ethernet) Address

MAC stands for Media Access Control

48 bit MAC address (for most LANs)

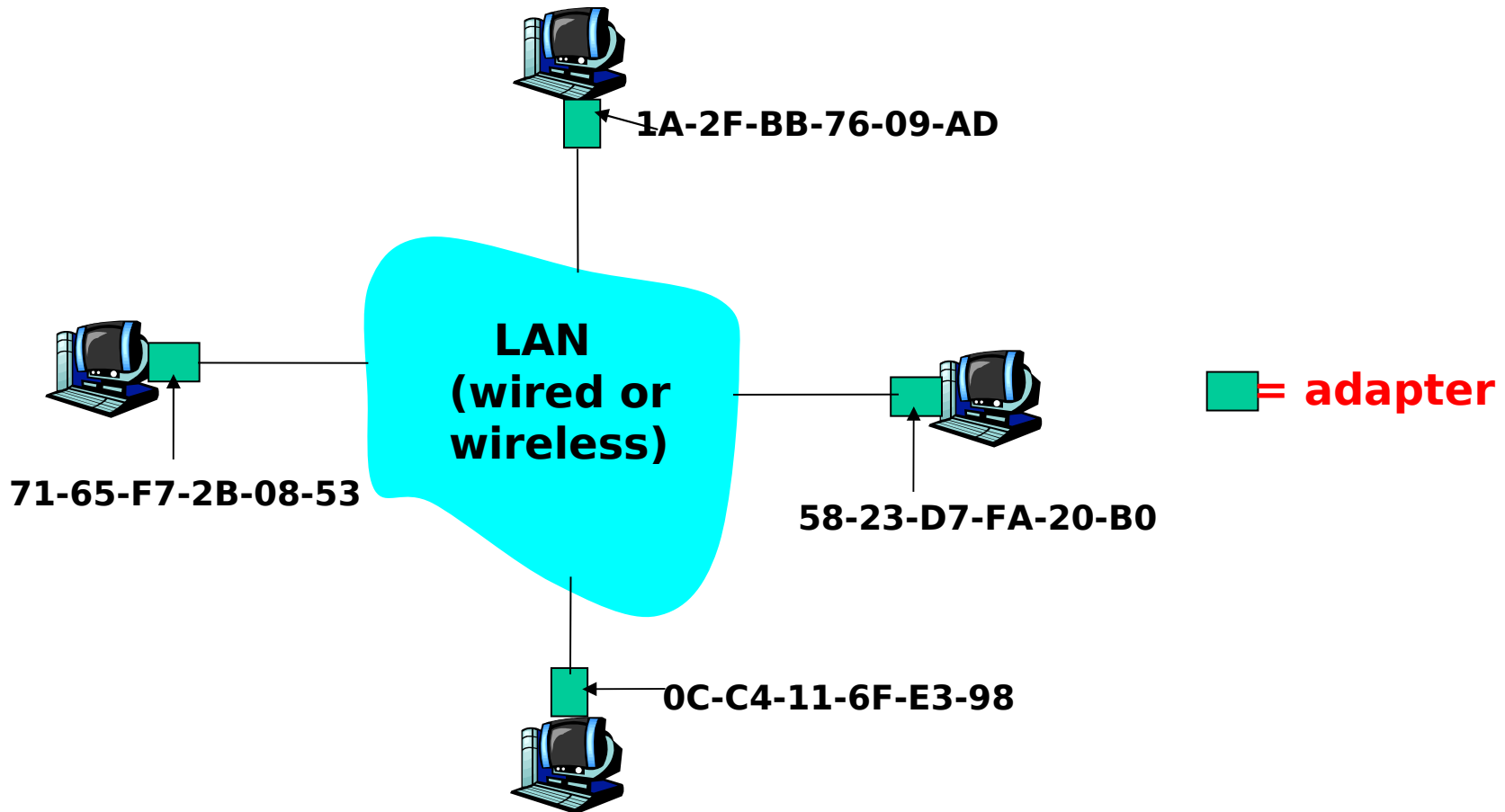
Burned in NIC ROM, also sometimes software settable

24 bits set for manufacturer, 24 bits for NIC adapter

Ex.: 00:E0:B8:9C:A6:60

MAC Addresses

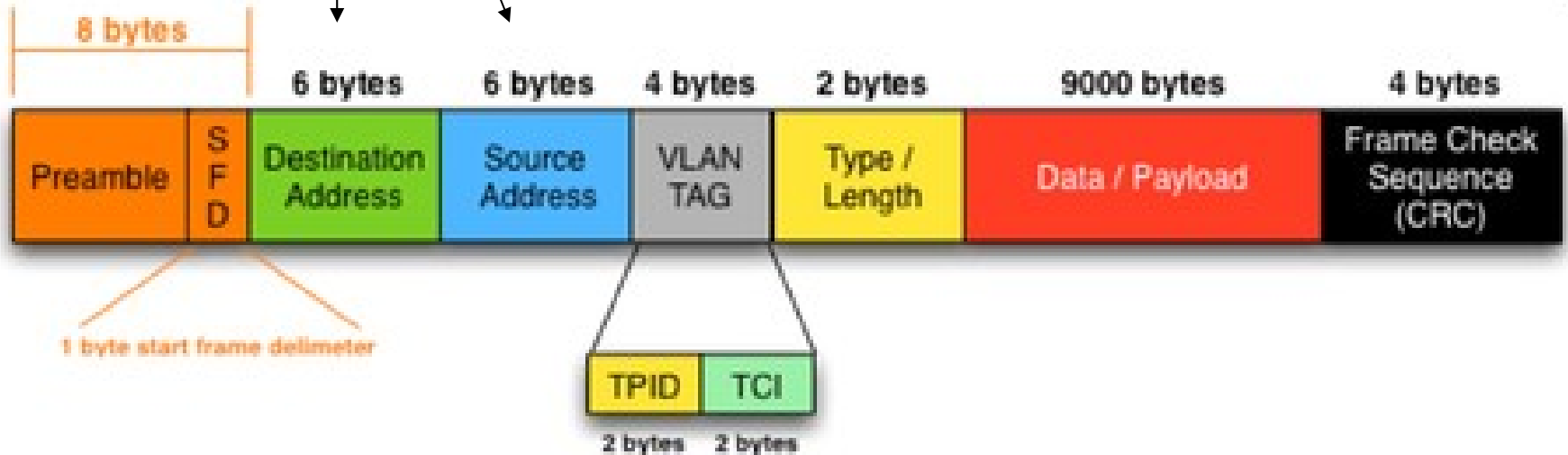
Each adapter on LAN has unique MAC address except for Broadcast address which is **FF-FF-FF-FF-FF-FF**



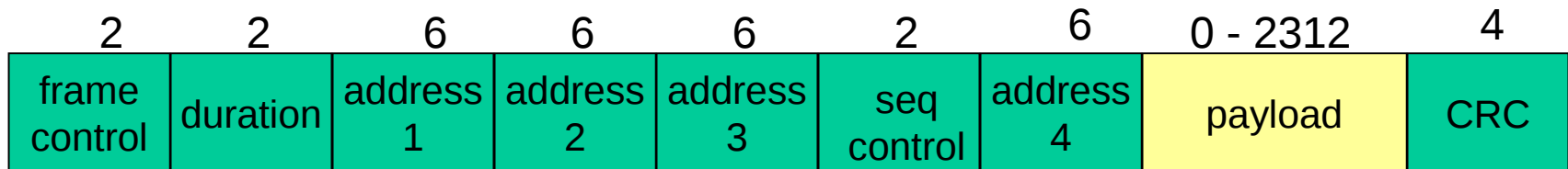
Ethernet Frame

Mac Addresses

9022 byte maximum Jumbo frame size with 802.1q VLAN Tag



802.11 Frame: Addressing



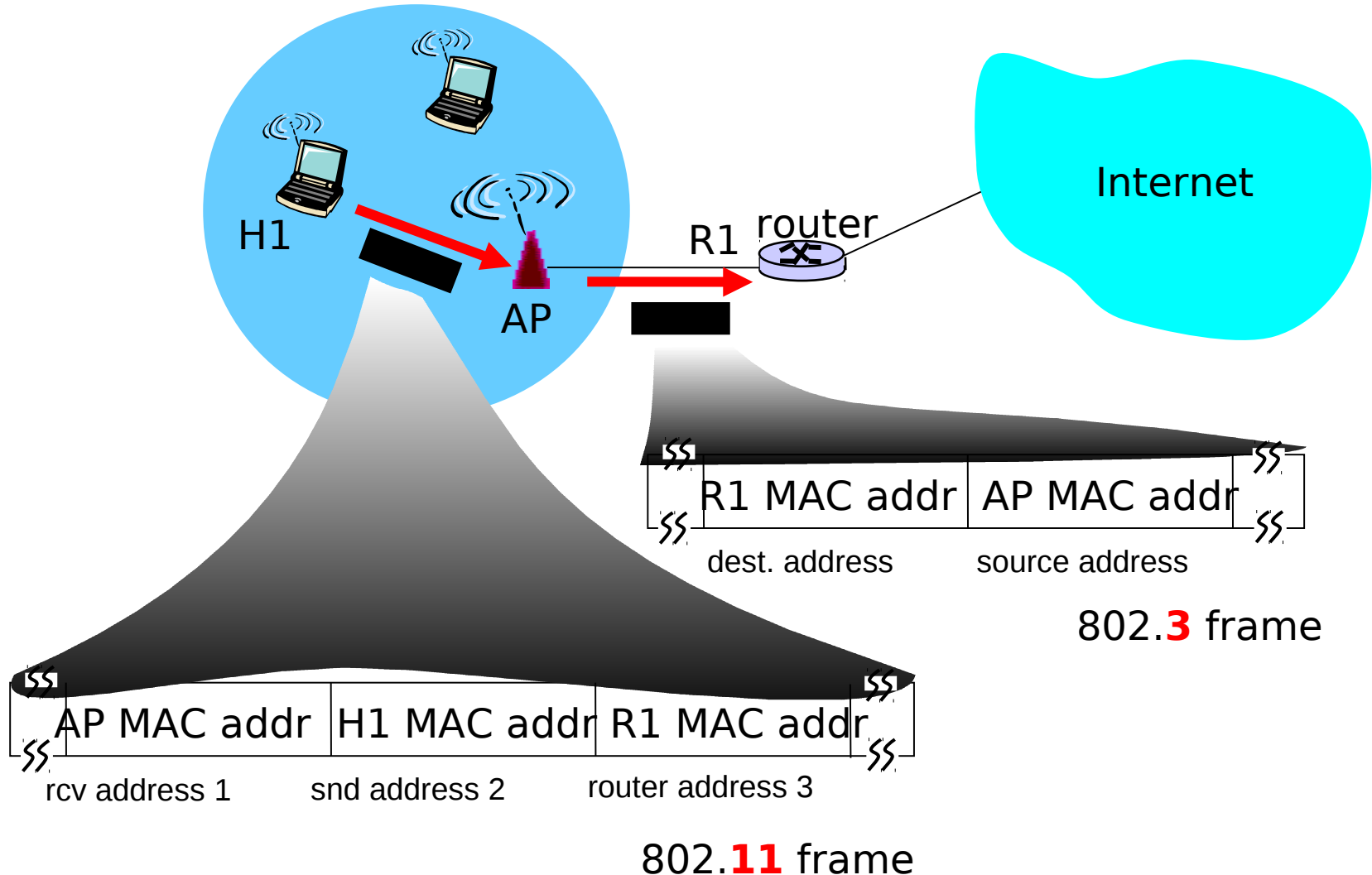
Address 1: MAC address of wireless host or AP to receive this frame

Address 2: MAC address of wireless host or AP transmitting this frame

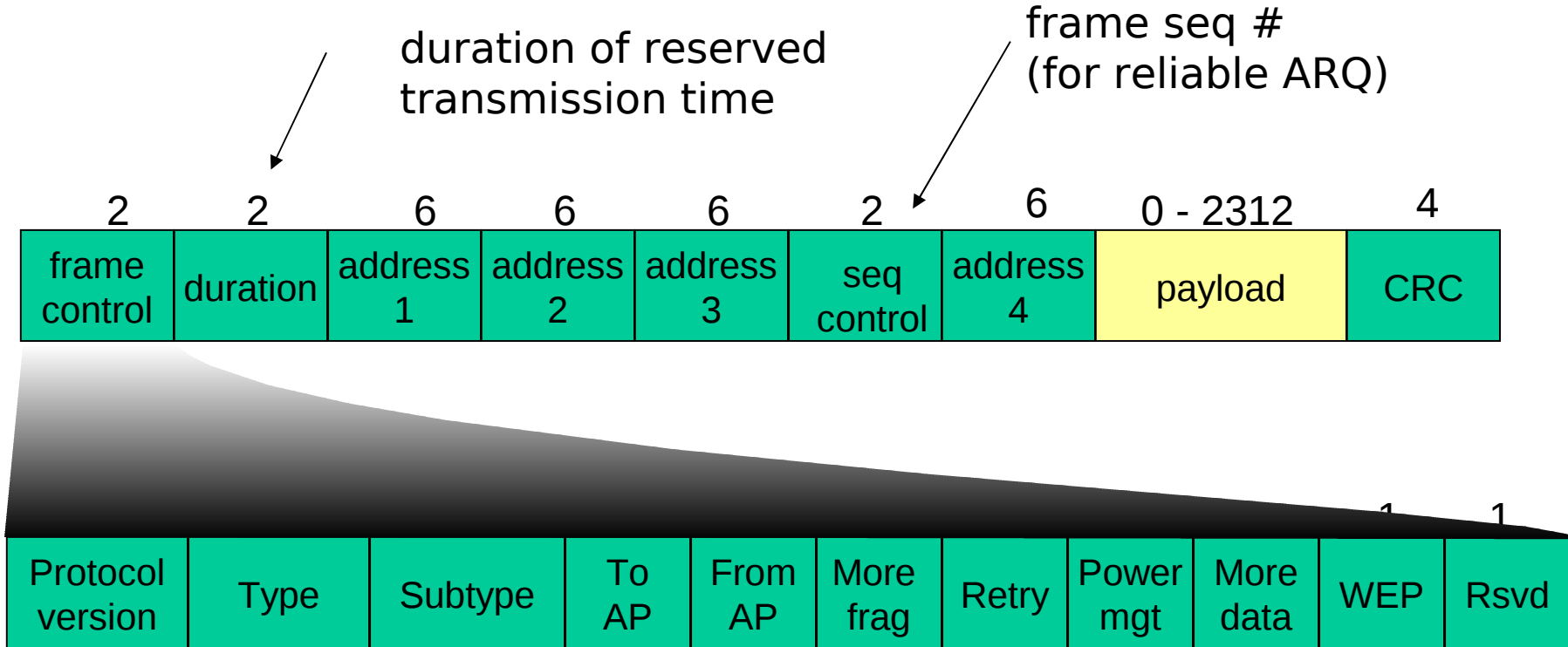
Address 3: MAC address of router interface to which AP is attached

Address 4: used only in ad hoc mode

802.11 frame: addressing



802.11 frame: more



2 bits 2 4 1 1 1 1 1 1 1 1

Protocol Version	Type	Subtype	To DS	From DS	More Fragments	Retry	Power Mgt.	More data	WEP	Order
------------------	------	---------	-------	---------	----------------	-------	------------	-----------	-----	-------

Protocol Version provides the current version of the 802.11 protocol used. Receiving STAs use this value to determine if the version of the protocol of the received frame is supported.

Type and Subtype determines the function of the frame.

There are three different frame type fields:

control, data, and management.

To DS and From DS indicates whether the frame is going to or exiting from the DS

This affects the order of the address fields see below

To DS field is 1 and From DS field is 0

Address 1 = BSSID

Address 2 = Source

Address 3 = Destination

To DS field is 0 and From DS field is 1

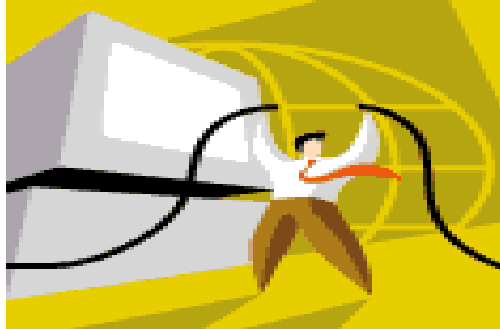
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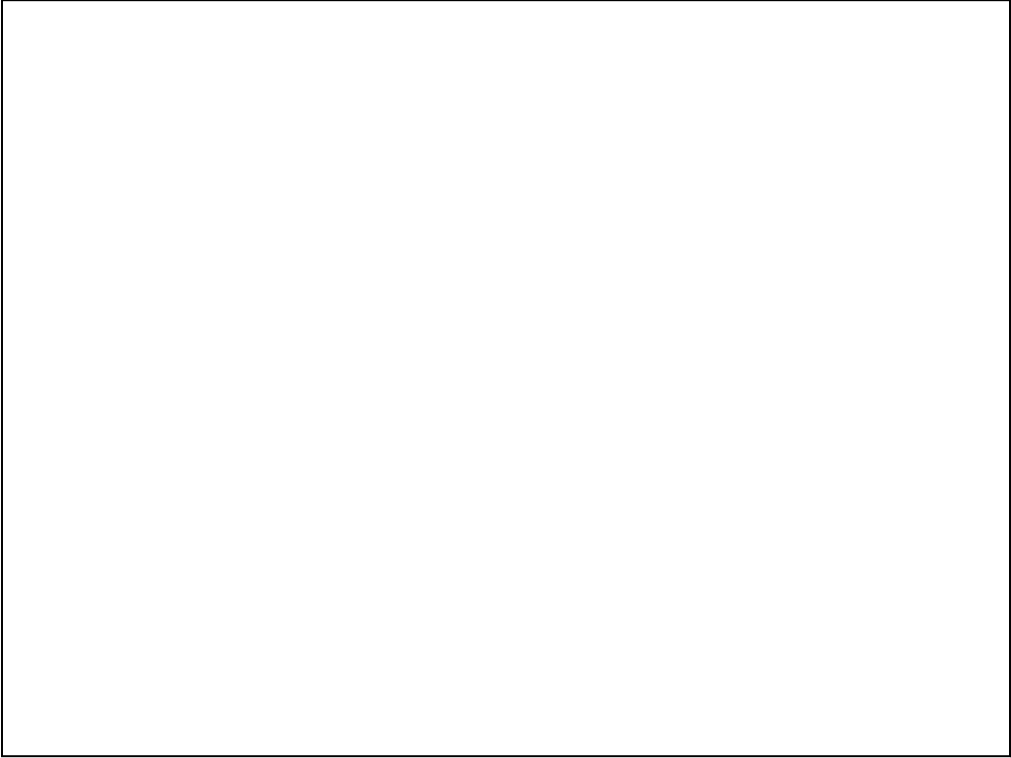
Address 3 = Source

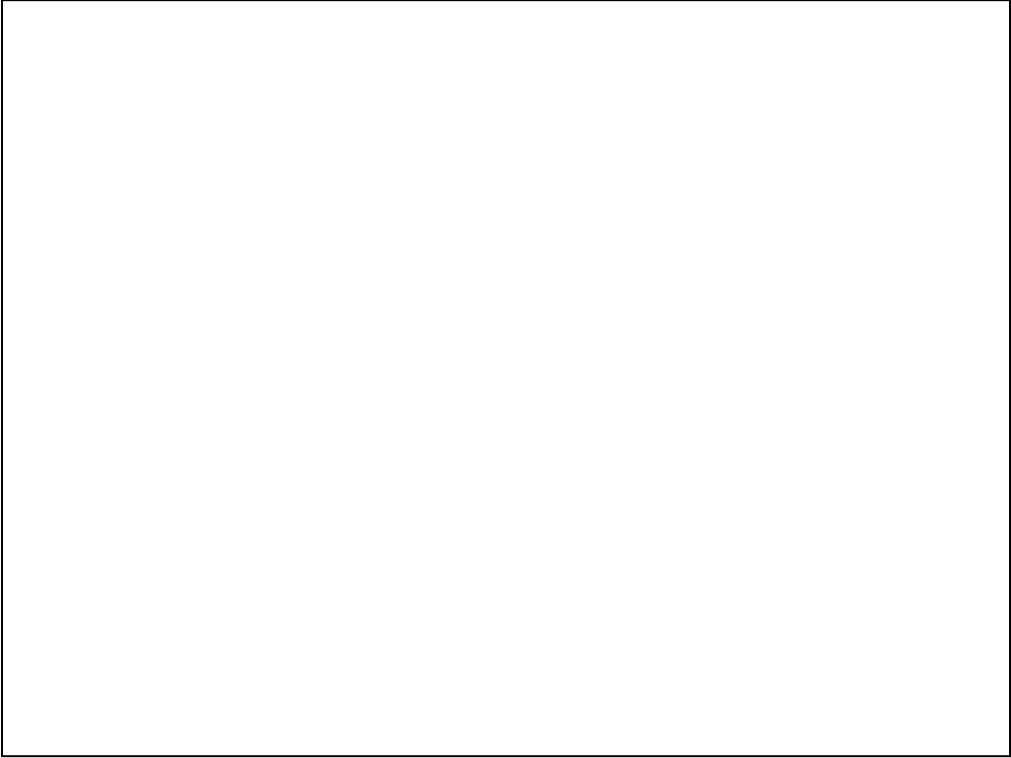
Summary

- 802.11 wireless more challenging because of disruptions to signal vs. wired
- However, mobility far outweighs the downside of interference and security
- No going back to wired when we can plug in during flights and have access to Facebook!
- Will see how wireless works in Wireshark lab

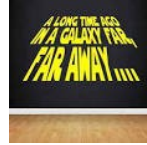


The End





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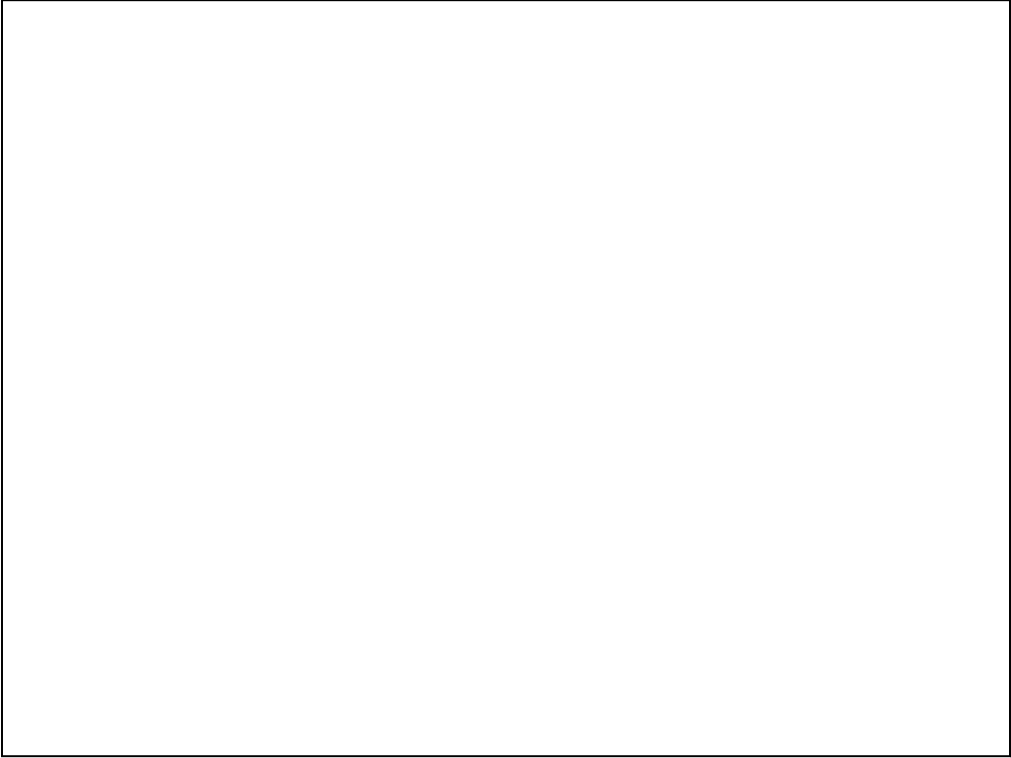
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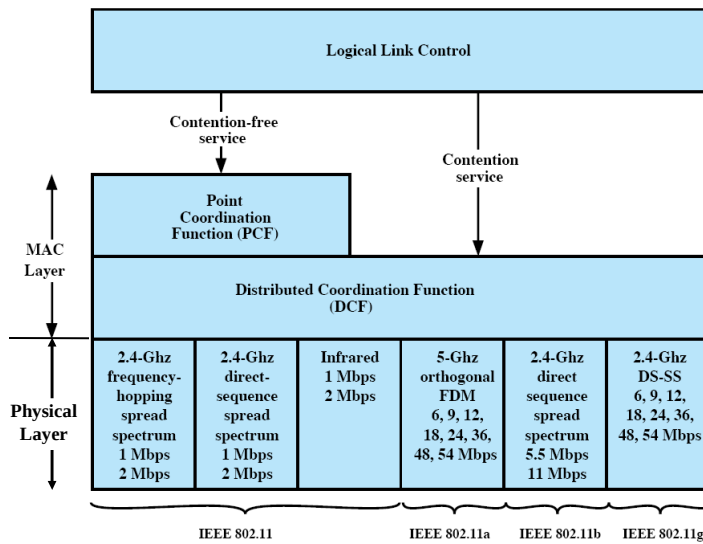
**IEEE 802.16™: Broadband Wireless Metropolitan
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802.11 Physical Layer

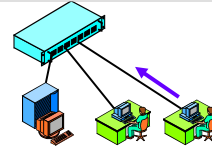
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IEEE 802.11 Protocol Architecture



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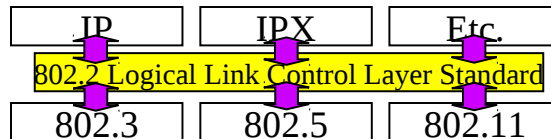


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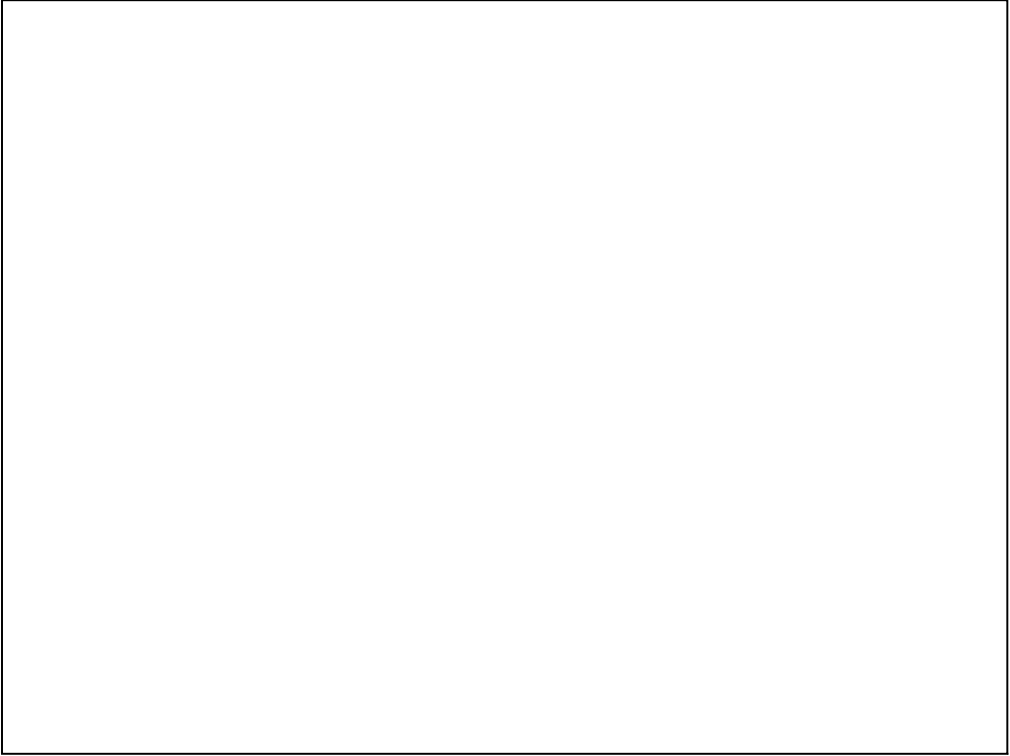
MAC Sublayer Functions

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Carrier Sense Multiple Access with Collision Avoidance

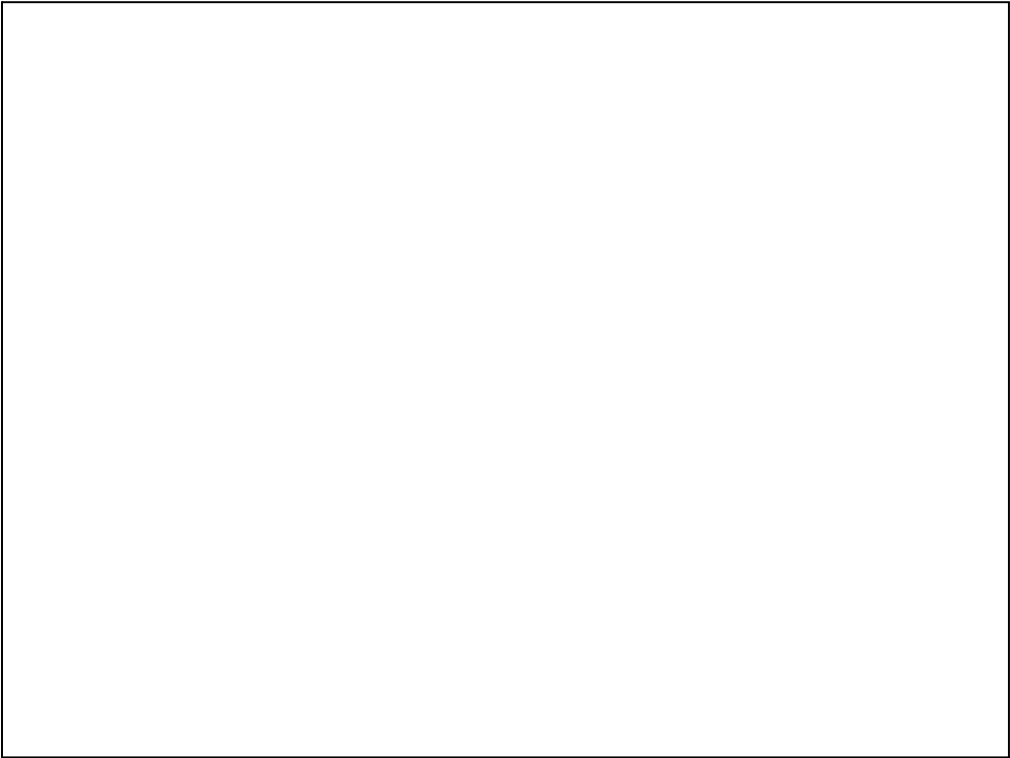
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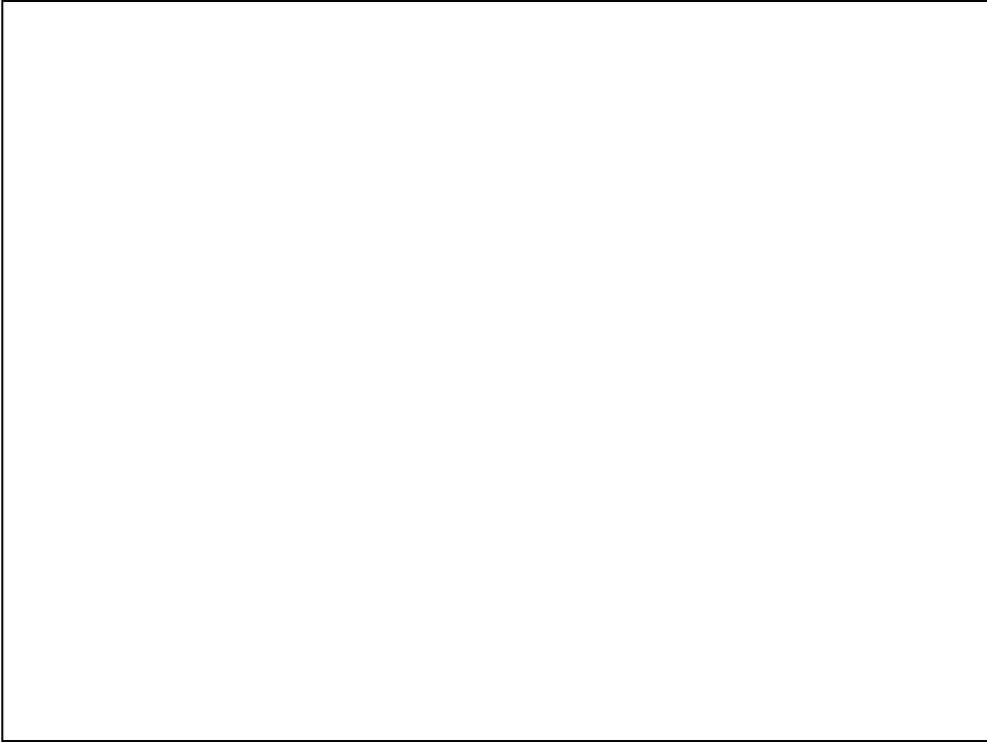
Recall that Ethernet uses CSMA/CD

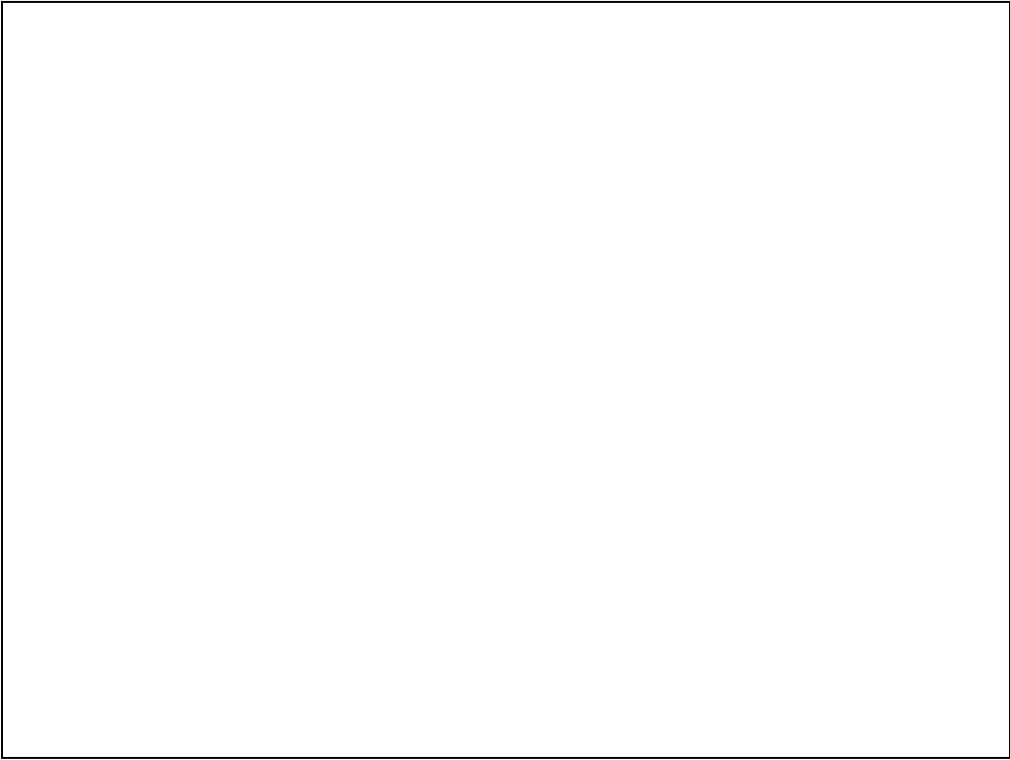


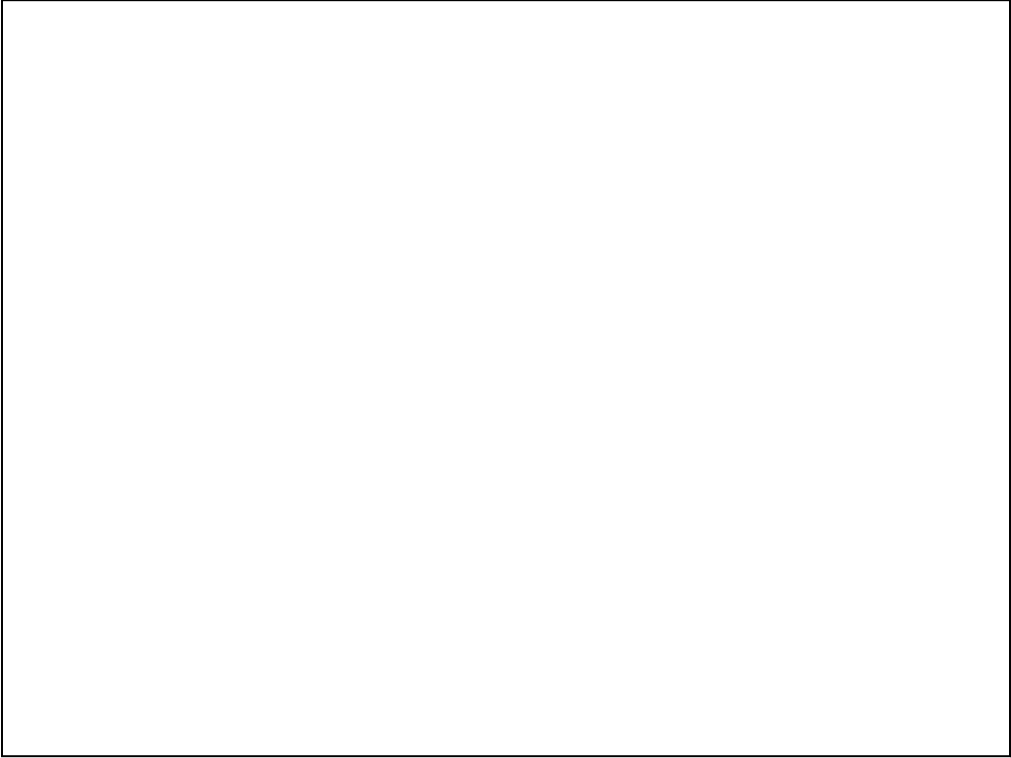
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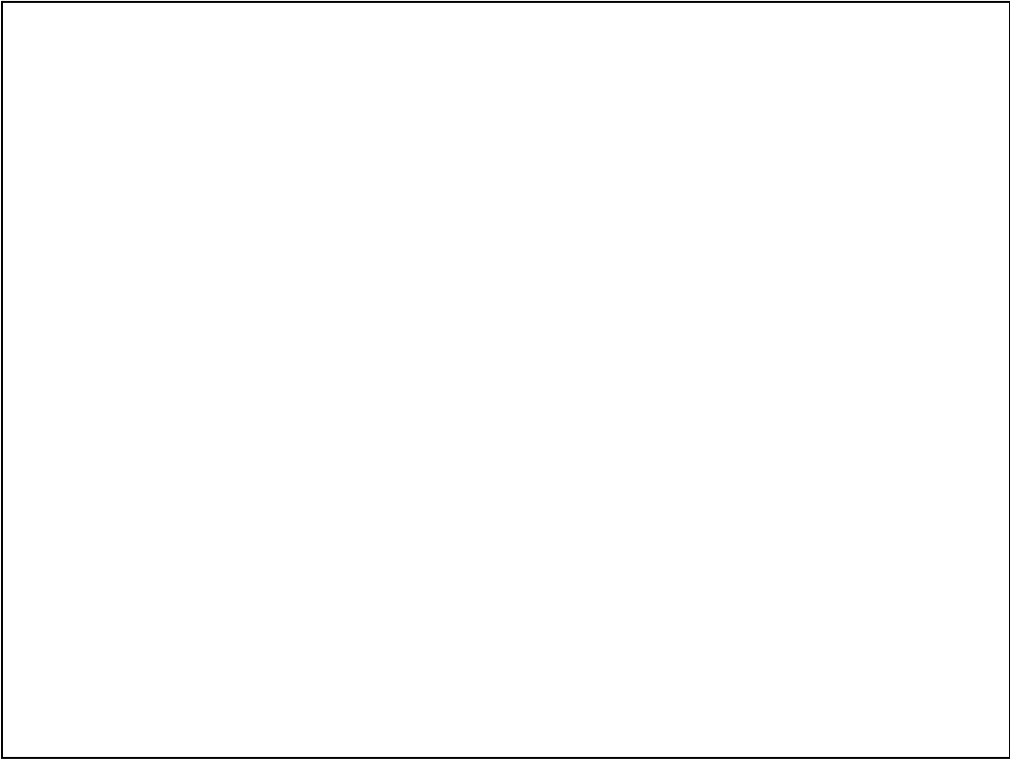
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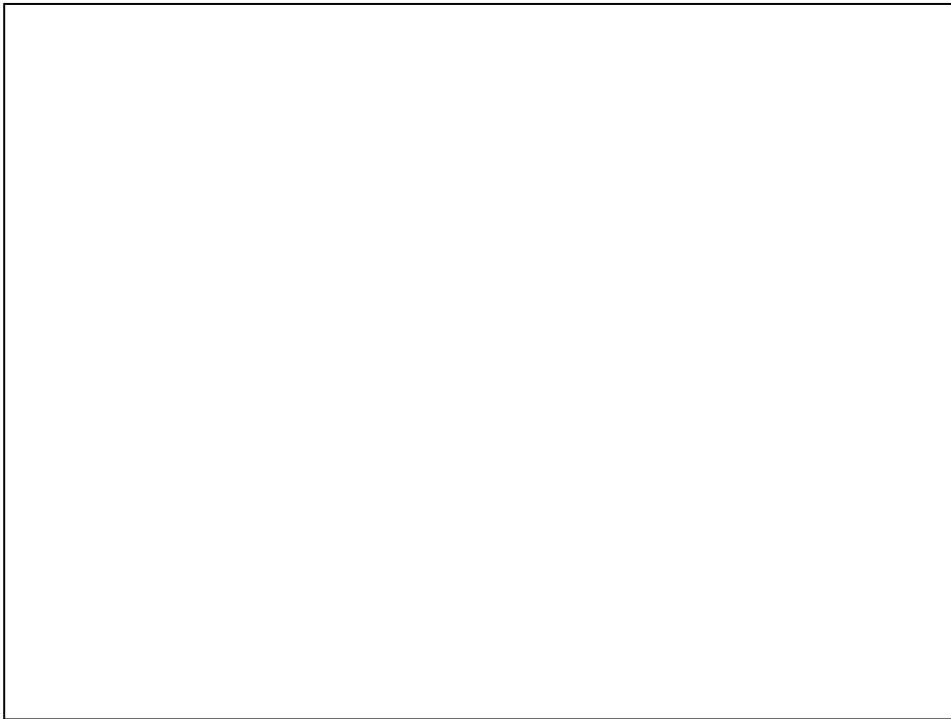
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- **Interference from other sources:** Standardized wireless network frequencies shared by other devices that can interfere
- **Multipath propagation:** Radio signal reflects off objects and ground, arriving at destination at slightly different times

.... make communication across link much more "difficult"



DIFS= Distributed interframe Space
PIFS=Point Coordination Interframe
Space
SIFS=Short Coordination Interframe
Space

802.11 Medium Access Control

- Shares the medium through coordination with other stations
 1. Sends control packets for coordination
Including Acks
 2. Sets and sends individual frame timers for all to see, each frame has its own timer
- Note: There is the 802.11 standard and then there is reality of did it actually get implemented!
 - Create a lot of details for a standard but not all of it is implemented

Wireless Collision Avoidance

STEPS

1. Have a frame to send
2. Wait a random time, until channel is idle
3. Sense it is idle for short time, called **DIFS** period
4. Sends frame, if gets through, destination waits a **SIFS** time and sends an **ACK**
5. Lack of an ACK back means frame failed
6. Sender then doubles backoff time, tries again
7. Continues until frame succeeds

Distributed Inter-frame Spacing (DIFS)
Short Inter-frame Spacing (SIFS)

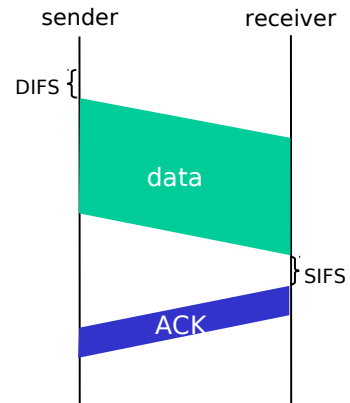
IEEE 802.11 MAC Protocol: CSMA/CA

802.11 sender

1. If sense channel idle for **DIFS** then
transmit entire frame
2. If sense channel busy then
 - a) start random backoff time
 - b) timer counts down
 - c) transmit when timer expires
 - d) if no ACK, increase random backoff interval, repeat 2 (frame failed)

802.11 receiver

- if frame received OK
return ACK after **SIFS** (ACK
needed due to hidden terminal problem)



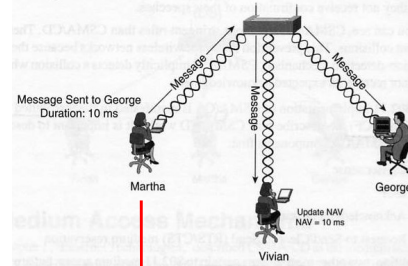
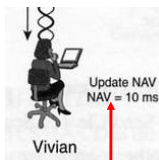
Distributed Inter-frame Spacing (DIFS)
Short Inter-frame Spacing (SIFS)

802.11 Frames have Timers

- **Network Allocation Vector (NAV) timer**
 - NAV is set when a frame sequence is sent
 - Says how long a sequence will take so other stations have an idea when the medium will be available
 - For example, a NAV for a data frame will also include the ACK back

Next Slides Demo this with Example ...

NAV Timer

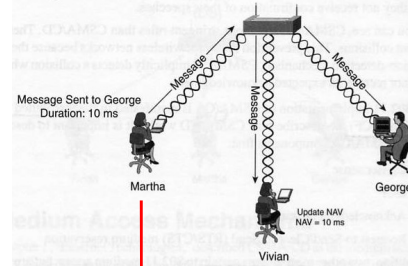
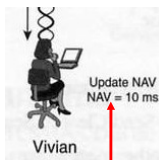


General 802.11 Frame (more on this later)

Frame Control	Duration/ID	Address 1	Address 2	Address 3	Sequence Control	Address 4	Frame Body	FCS
2 Bytes	2 Bytes	6 Bytes	6 Bytes	6 Bytes	2 Bytes	6 Bytes	0 - 2312 Bytes	4 Bytes

- All stations have a **NAV (Network Allocation Vector) timer**
Protects frames from interruption
- **Example: Martha sends a frame to George**
- Since wireless medium is “broadcast-based” shared medium, all stations including Vivian receive frame
- Vivian updates her NAV timer with duration value
- Vivian will not attempt to transmit until her NAV is decremented to 0.
- Stations will only update their NAV when duration field value received is greater than their current NAV

Duration Field



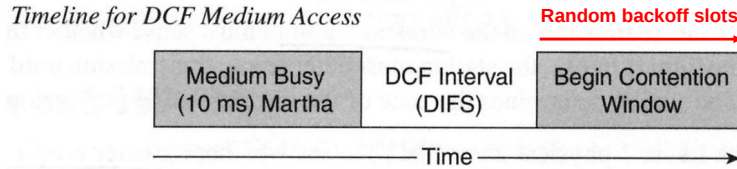
General 802.11 Frame (more on this later)

Frame Control	Duration/ID	Address 1	Address 2	Address 3	Sequence Control	Address 4	Frame Body	FCS
2 Bytes	2 Bytes	6 Bytes	6 Bytes	6 Bytes	2 Bytes	6 Bytes	0 - 2312 Bytes	4 Bytes

- **Duration/ID field** – Number of microseconds medium is expected to remain busy for transmission currently in progress
 - **Transmitting device sets Duration time in microseconds**
 - **Includes time to:**
 - **Transmit this frame to AP (or to the client if an AP)**
 - **Includes returning ACK**
- All stations monitor this field!
- All stations update their **NAV** (Network Allocation Vector) timer

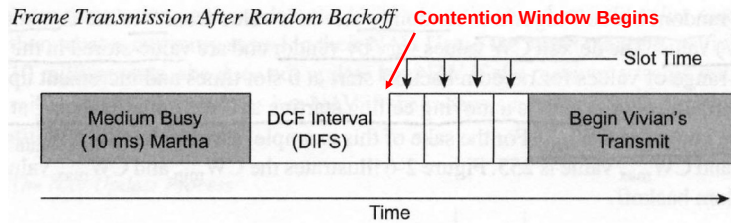
Wanting to transmit (1/3)

Timeline for DCF Medium Access



- Station wanting to transmit.
- **Carrier Sensing**
 - **Physical: Physically senses medium is idle**
 - **Virtual: NAV timer is 0**
- Waits **DIFS** (Distributed or DCF Interframe Space)
 - **Minimum amount of medium idle time until contention-based services begin.**
 - **Once DIFS is over, stations can contend for access.**
- **Contention window** begins.
 - **Uses random backoff algorithm to determine when it can attempt to access the medium. (next)**

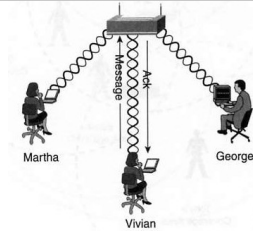
Wanting to transmit (2/3)



- The **random backoff algorithm** randomly selects a value from **0 to 255**. Maximum value varies by vendor.
- The random value is number of **802.11 slot times** the station must wait after the DIFS, during the contention window before it may transmit.
- **Stations pick a random slot** and wait for that slot before attempting to access the medium.
- With several stations attempting to transmit, the station that picks the **lowest slot, lowest random number, wins**.

Wanting to transmit (3/3)

Others
update NAV



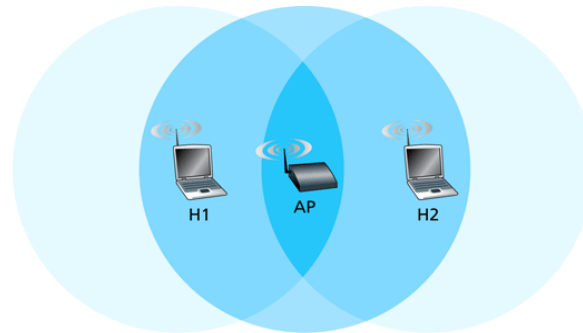
General 802.11 Frame (more on this later)

Frame Control	Duration/ID	Address 1	Address 2	Address 3	Sequence Control	Address 4	Frame Body	FCS
2 Bytes	2 Bytes	6 Bytes	6 Bytes	6 Bytes	2 Bytes	6 Bytes	0 - 2312 Bytes	4 Bytes

- **Station transmits**, setting the **Duration ID** to the time needed to transmit data, ACK and IFSSs.
- **Other stations** with higher slots will see the new transmission and wait to transmit.
- If **frame arrives at AP** (assuming the transmitter is a station), then an **ACK will be returned**
- If there is **not an ACK received**, the sending station assumes there has been a collision
 - **If two stations have the same lowest slot time and both transmit, then a collision occurs**
- Stations will **update its retry counter** (double) to determine a **new randomly selected slot time** and **process starts all over again**

Hidden Terminal Problem in WLANs

- Both H1 and H2 transmit at same time
- Signals collide at AP, H1 can't detect H2



Collision is the darker blue

Figure 6.11 ♦ Hidden terminal example: H1 is hidden from H2, and vice versa.

Avoiding collisions: RTS/CTS

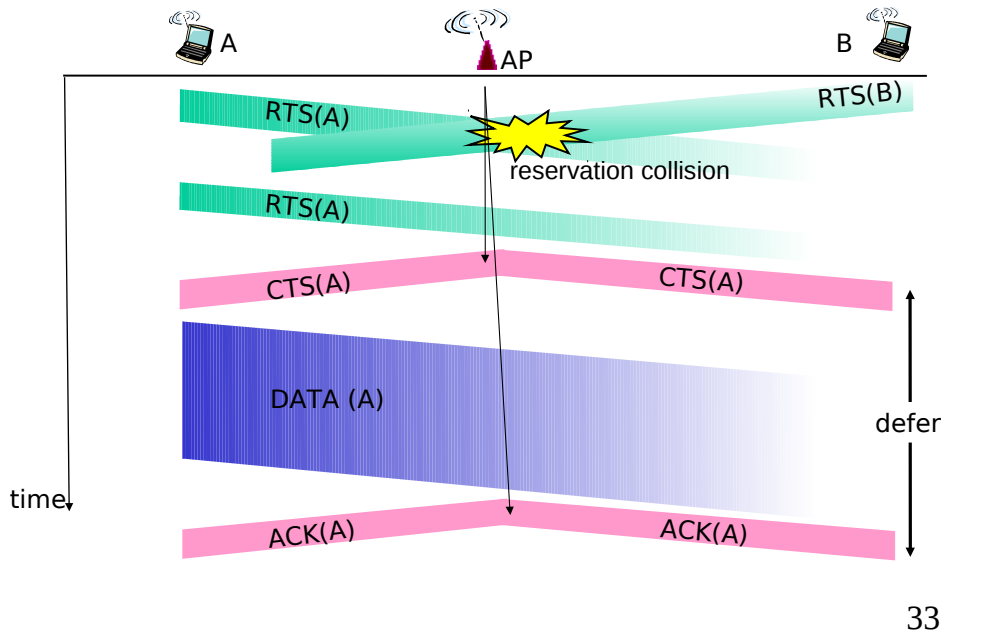
Another Idea: Allow sender to “reserve” channel rather than random access of data frames: avoid collisions of long data frames

Sender first transmits *small* request-to-send (RTS) packets using CSMA, say want to transmit to AP

- RTSs may still collide with each other (but they’re short)
- AP broadcasts clear-to-send (CTS) in response to RTS
- RTS heard by all nodes
 - Sender transmits data frame
 - Other stations defer transmissions

Avoids data frame collisions completely
using small reservation packets!

Collision Avoidance: RTS-CTS exchange

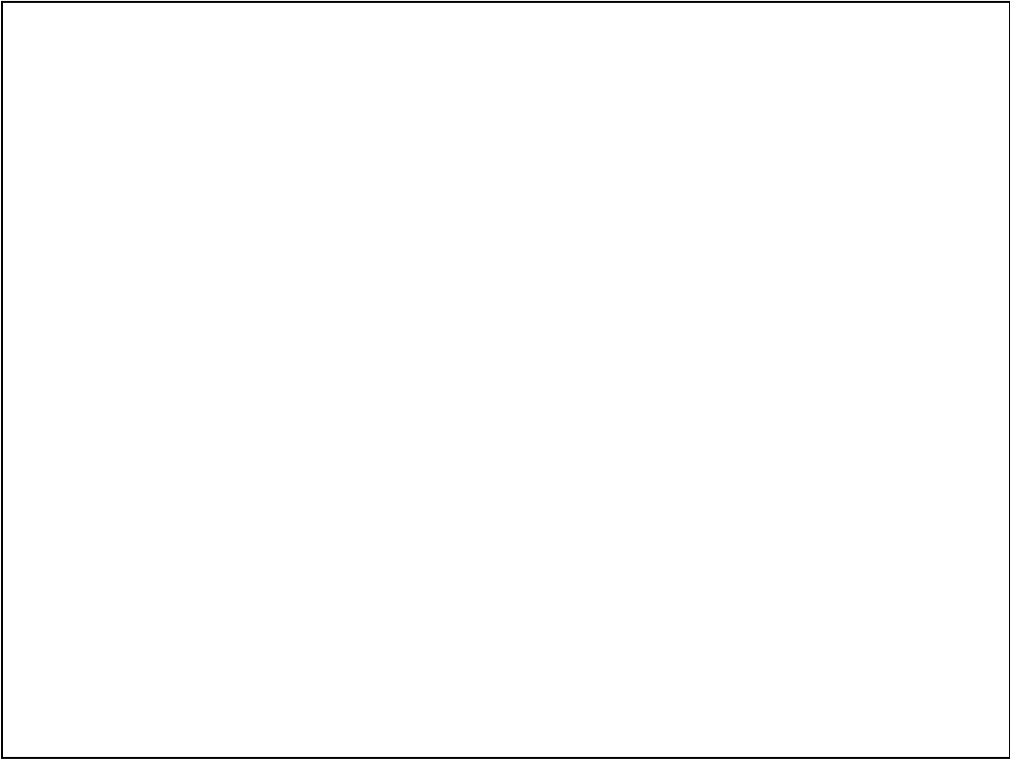


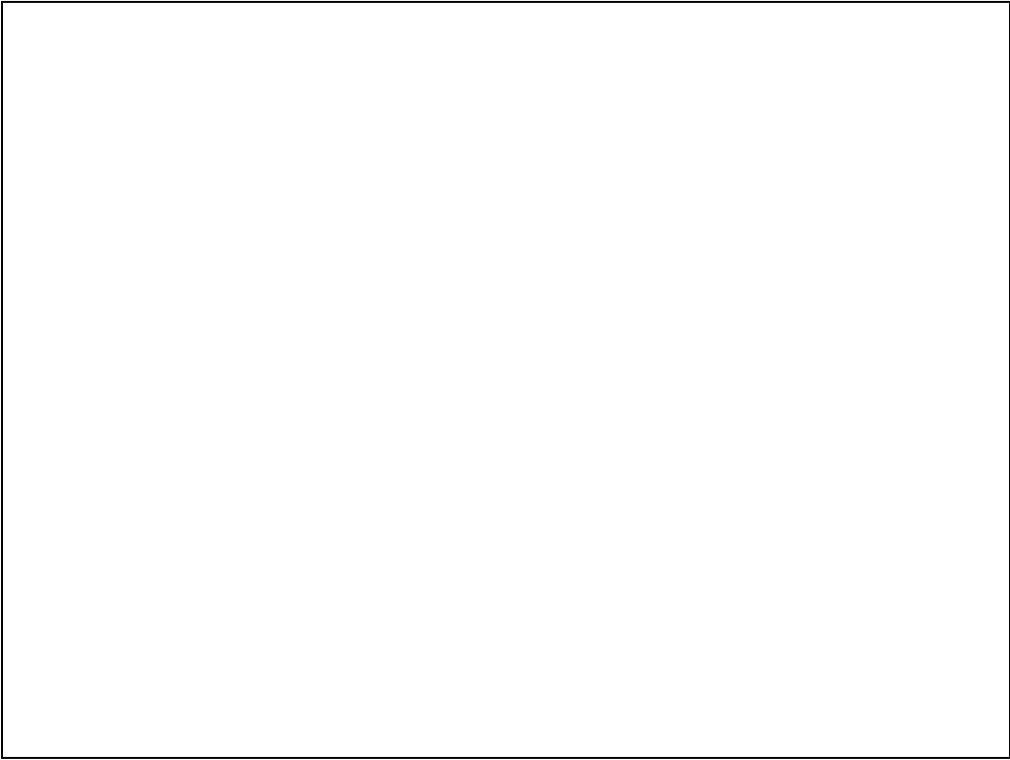
RTS/CTS in practice

- 802.11 standardized both CSMA/CA and RTS/CTS
- In practice, most operators **disable** RTS/CTS
 - **Very high overhead!**
 - RTS/CTS packets sent at “base rate” (often 1Mbit)
 - Neighboring AP's are often configured to use non-overlapping channels, so hidden terminals on downlink are rare

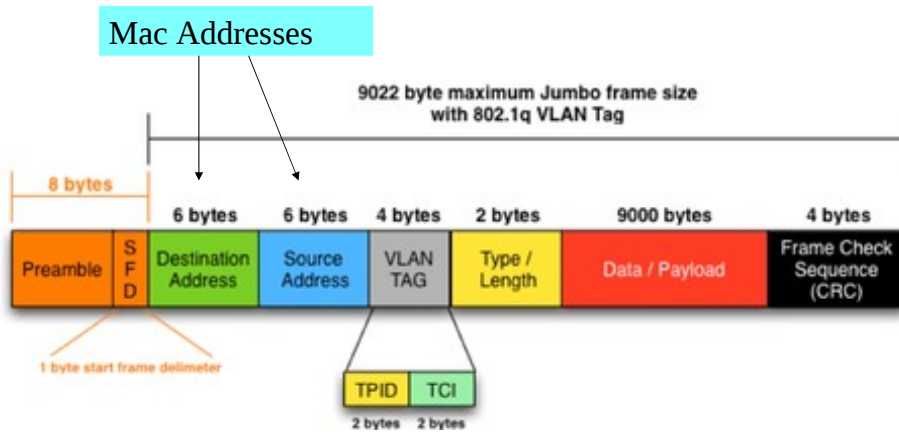


MAC Addresses Review

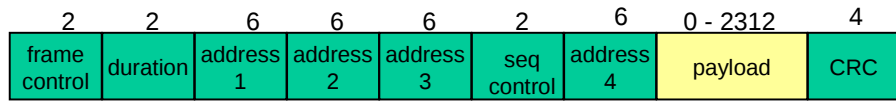




Ethernet Frame



802.11 Frame: Addressing



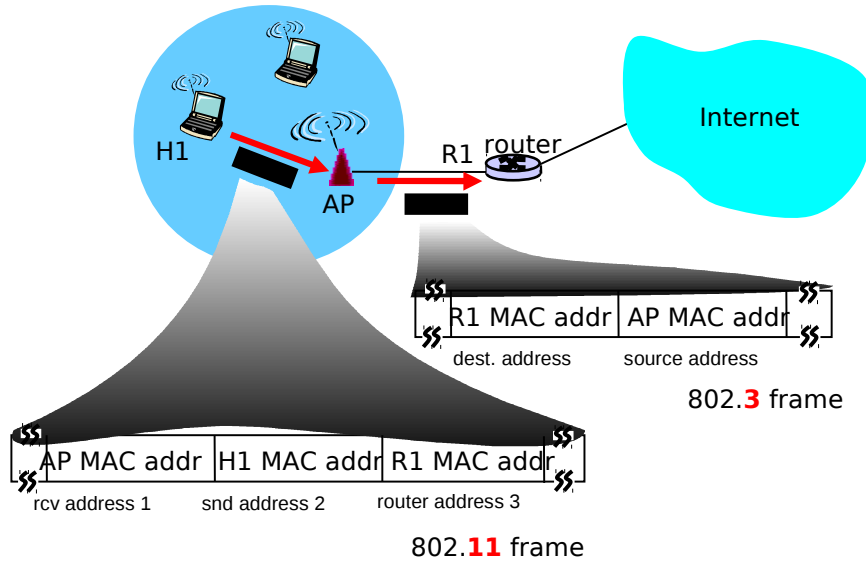
Address 1: MAC address of wireless host or AP to receive this frame

Address 2: MAC address of wireless host or AP transmitting this frame

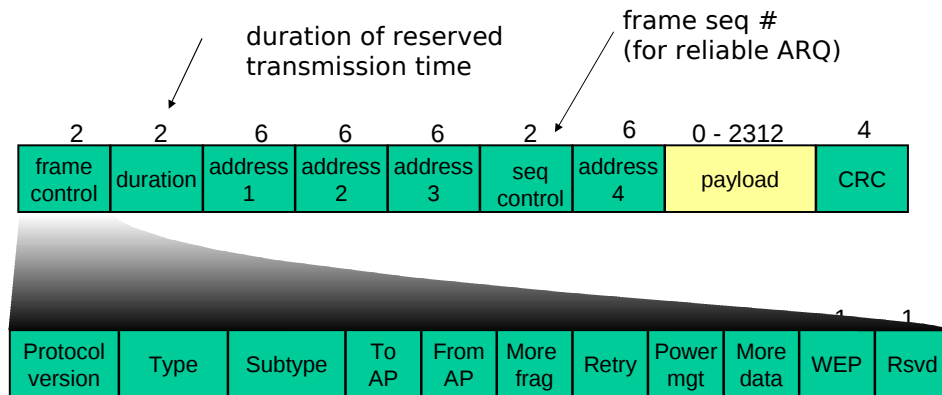
Address 3: MAC address of router interface to which AP is attached

Address 4: used only in ad hoc mode

802.11 frame: addressing



802.11 frame: more



2 bits	2	4	1	1	1	1	1	1	1	1
Protocol Version	Type	Subtype	To DS	From DS	More Fragments	Retry	Power Mgt.	More data	WEP	Order

Protocol Version provides the current version of the 802.11 protocol used. Receiving STAs use this value to determine if the version of the protocol of the received frame is supported.

Type and Subtype determines the function of the frame.

There are three different frame type fields:

control, data, and management.

To DS and From DS indicates whether the frame is going to or exiting from the DS

This affects the order of the address fields see below

To DS field is 1 and From DS field is 0

Address 1 = BSSID

Address 2 = Source

Address 3 = Destination

To DS field is 0 and From DS field is 1

Address 1 = Destination

Address 2 = BSSID

Address 3 = Source

